

Inverse theorems and approximate structure

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Question

What do $\begin{cases} 99\% \\ 1\% \end{cases}$ homomorphisms look like?

The structure of 99% homomorphisms

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Take a true homomorphism $\phi: G_1 \rightarrow G_2$ and modify $(\varepsilon/3)|G_1|$ of the values $\phi(x)$.

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If

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then there exists a homomorphism $\phi: G_1 \rightarrow G_2$ such that

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Application

Property testing.

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Is it true that there exists a homomorphism $\phi: G_1 \rightarrow G_2$ with $\mathbb{P}_{x \in G_1} [f(x) = \phi(x)] \geq F(\delta)$?

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Consider $f: \mathbb{Z} \rightarrow \mathbb{Z}$,

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We have

$\mathbb{P}[\lfloor \sqrt{2}(x+y) \rfloor = \lfloor \sqrt{2}x \rfloor + \lfloor \sqrt{2}f(y) \rfloor] \approx 1/2$
(depending on whether the carry bit is 0 or 1).

So no.

The structure of 1% homomorphisms (contd.)

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Theorem (Samorodnitsky 2007)

If $G_1 = \mathbb{F}_2^n$, $G_2 = \mathbb{F}_2^m$ then yes, $\mathbb{P}_{x \in G_1} [f(x) = \phi(x)] \geq F(\delta)$ for some homomorphism ϕ .

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Conjecture (Ruzsa; Conjecture A)

You can take $F(\delta) = \delta^C$ for some constant $C > 0$.

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Weak (statistical structure) implies *strong* (algebraic) structure.

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Let G be an abelian group and $A \subseteq G$ a finite subset.

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What do $\begin{cases} 99\% \\ 1\% \end{cases}$ subgroups look like in general?

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Proposition (Freiman (99%))

If $|A + A| < (1 + \varepsilon)|A|$ and $\varepsilon < 1/2$ then there is a subgroup $H \leq G$ and $a \in G$ such that $A \subseteq a + H$ and $|H| \leq (1 + \varepsilon)|A|$.

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Theorem (Ruzsa (1%))

If $A \subseteq \mathbb{F}_2^n$ and $|A + A| \leq K|A|$ then there exists a subgroup $H \leq \mathbb{F}_2^n$ and $a \in \mathbb{F}_2^n$ such that $A \subseteq a + H$ and $|H| \leq K^2 2^{K^4} |A|$.

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Conjecture (Marton; the “polynomial Freiman–Ruzsa conjecture”; Conjecture B)

If $A \subseteq \mathbb{F}_2^n$ and $|A + A| \leq K|A|$ then there exists a subgroup $H \leq \mathbb{F}_2^n$ and $a_1, \dots, a_m \in \mathbb{F}_2^n$ such that

$$A \subseteq \bigcup_{i=1}^m (a_i + H)$$

and $|H| \leq |A|$ and $m \leq 2K^C$.

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Idea for \Leftarrow .

Given $f: \mathbb{F}_2^n \rightarrow \mathbb{F}_2^m$ consider its *graph*

$$\Gamma = \{(x, f(x)) : x \in \mathbb{F}_2^n\} \subseteq \mathbb{F}_2^n \times \mathbb{F}_2^m.$$

Then Γ is a (probabilist's) 1% subgroup. □

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Theorem (Gowers, Green, M., Tao 2024+)

Conjecture B is true (with $C = 12$).

Very brief outline of proof

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 $|A|$

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Yet another notion of approximate structure:

$$\begin{array}{l|l} \text{Set } A \subseteq G & \text{Random variable } X \text{ on } G \\ |A| & \text{Shannon entropy } H[X] \\ |A + A| \leq K|A| & H[X_1 + X_2] - \frac{1}{2}H[X_1] - \frac{1}{2}H[X_2] \leq k \end{array}$$

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Recursive proof: given X_1, X_2 , either $k = 0$ or find X'_1, X'_2 with $k(X'_1, X'_2) \leq 0.99k(X_1, X_2)$.

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Definition

For G_1, G_2 abelian groups, call $f: G_1 \rightarrow G_2$ a *polynomial of degree $\leq s$* if $\partial_{h_1} \dots \partial_{h_{s+1}} f(x) = 0 \forall x, h_1, \dots, h_{s+1} \in G_1$.

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$f: G_1 \rightarrow G_2$ is a *probabilist's approximate polynomial* if

$$\mathbb{P}_{x, h_1, \dots, h_{s+1} \in G_1} [\partial_{h_1} \dots \partial_{h_{s+1}} f(x) = 0] \geq \delta.$$

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$g: G \rightarrow \mathbb{C}$, $\|g\|_\infty \leq 1$ (e.g., $g(x) = e^{2\pi i f(x)}$) is an *analyst's approximate polynomial* if

$$\mathbb{E}_{x, h_1, \dots, h_{s+1}} \Delta_{h_1} \dots \Delta_{h_{s+1}} g(x) \geq \delta$$

where $\Delta_h g(x) = g(x) \overline{g(x+h)}$

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$$\|g\|_{U^{s+1}}^{2^{s+1}} := \mathbb{E}_{x, h_1, \dots, h_{s+1}} \Delta_{h_1} \dots \Delta_{h_{s+1}} g(x) \geq \delta$$

where $\Delta_h g(x) = g(x) \overline{g(x+h)}$ and $\|g\|_{U^{s+1}}$ is called a *Gowers uniformity norm*.

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$$\mathbb{E}_{x, h_1, \dots, h_{s+1} \in G} \Delta_{h_1} \dots \Delta_{h_{s+1}} g(x) \geq \delta \quad (\dagger)$$

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If $g: G \rightarrow \mathbb{C}$, $\|g\|_\infty \leq 1$, obeys (\dagger) , \exists a true polynomial $\phi: G \rightarrow \mathbb{R}/\mathbb{Z}$,

$$|\mathbb{E}_{x \in G} g(x) e^{-2\pi i \phi(x)}| \geq F(\delta)?$$

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Example ($s = 1$)

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Higher order Fourier analysis

$$\mathbb{E}_{x, h_1, \dots, h_{s+1} \in G} \Delta_{h_1} \dots \Delta_{h_{s+1}} g(x) \geq \delta \quad (\dagger)$$

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If $g: G \rightarrow \mathbb{C}$, $\|g\|_\infty \leq 1$, obeys (\dagger) , \exists a true polynomial $\phi: G \rightarrow \mathbb{R}/\mathbb{Z}$,

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The case $s > 1$ is called “higher-order Fourier analysis”.

The inverse theorem(s)

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Theorem (Tao–Ziegler (2011);)

If $g: \mathbb{F}_p^n \rightarrow \mathbb{C}$, $\|g\|_\infty \leq 1$ and (\dagger) holds, \exists a degree $\leq s$ polynomial $\phi: G \rightarrow \mathbb{R}/\mathbb{Z}$,

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If $g: \mathbb{Z}/n\mathbb{Z} \rightarrow \mathbb{C}$, $\|g\|_\infty \leq 1$ and (\dagger) holds, then there exists a “nilsequence” $\psi: G \rightarrow \mathbb{C}$ such that

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The inverse theorem and approximate polynomials

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In particular, the structure theorem for 1% linear functions

\Rightarrow the inverse theorem for quadratic Fourier analysis (i.e., the U^3 -norm).

Question (Green–Tao)

Do there exist $x, a \in \mathbb{Z}$ such that $x, x + a, x + 2a, x + 3a$ are all prime?

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[END]