

Kazhdan's property (T) for $\text{Aut}(\mathbf{F}_n)$ and $\text{EL}_n(\mathcal{R})$

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Harvard Math Picture Language Seminar, 2022.10.11

C^* -algebras are an esoteric subject — “the most abstract nonsense that exists in mathematics,” in Casazza’s words. “Nobody outside the area knows much about it.”

Quanta Magazine: *‘Outsiders’ Crack 50-Year-Old Math Problem.*

<http://www.quantamagazine.org/>

computer-scientists-solve-kadison-singer-problem-20151124

Kazhdan's property (T)

Theorem (Kazhdan 1967)

Any simple Lie group G of real rank ≥ 2 (e.g., $G = \mathrm{SL}_n(\mathbb{R})$, $n \geq 3$) and its lattice Γ (e.g., $\Gamma = \mathrm{SL}_n(\mathbb{Z})$, $n \geq 3$) have **property (T)**.

$\rightsquigarrow \Gamma$ is finitely generated and has finite abelianization.

Definition (for a discrete group Γ)

Γ has (T) $\stackrel{\text{def}}{\iff} \exists S \subset \Gamma$ finite $\exists \kappa > 0$ s.t. $\forall (\pi, \mathcal{H})$ unitary rep'n and $\forall v \in \mathcal{H}$

$$d(v, \mathcal{H}^\Gamma) \leq \kappa^{-1} \max_{s \in S} \|v - \pi(s)v\|.$$

$\iff \Gamma$ is f.g. & $\forall S \subset \Gamma$ generating $\exists \kappa = \kappa(\Gamma, S) > 0$ s.t. \dots

The optimal $\kappa(\Gamma, S)$ is called the **Kazhdan constant** for (Γ, S) .

- Property (T) inherits to finite-index subgroups and quotient groups.
- \mathbb{Z} (or any infinite amenable group) does not have property (T).
 $\because \frac{1}{\sqrt{2k+1}} 1_{[-k,k]} \in \ell^2(\mathbb{Z})$ is asymp. \mathbb{Z} -invariant, but $\ell^2(\mathbb{Z})^\mathbb{Z} = \{0\}$.

\rightsquigarrow Any f.i. subgroup with property (T) has finite abelianization.

An application of property (T): Expander graphs

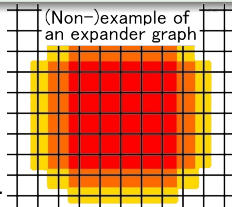
Explicit construction of expanders (Margulis 1973)

$\Gamma = \langle S \rangle$, X a finite set, and $\Gamma \curvearrowright X$ transitively

\rightsquigarrow Schreier graph: Vertices = X and Edges = $\{\{x, sx\} : x \in X, s \in S\}$
is a ε -**expander** for $\varepsilon = \frac{\kappa(\Gamma, S)^2}{2}$. Namely, for $\forall A \subset X$ one has

$$|\partial A| \geq \varepsilon |A| (1 - \frac{|A|}{|X|}).$$

- For $\mathcal{N}_k(A) := \{x \in X : d(x, A) \leq k\}$,
 $|\mathcal{N}_k(A)| \geq (1 + \varepsilon)^k |A|$ until it reaches $\frac{1}{2}|X|$.
After that $|\mathcal{N}_k(A)^c|$ decreases by a factor $1 + \varepsilon$.
- Random walk on X has mixing time $O(\log |X|)$.
- Want a large ε -expander with fixed degree and ε .



Expander Sampling Theorem (Gillman 1998):

For $\forall f \in \ell_\infty(X)$ with $\|f\|_\infty \leq 1$ and $m = \frac{1}{|X|} \sum f(x)$,

Simple Random Walk x_0, x_1, \dots on X satisfies

$$\mathbb{P}(|\frac{1}{T} \sum_{k=0}^{T-1} f(x_k) - m| > \delta) < \exp(-C_{d,\varepsilon} \delta^2 T).$$

Some examples of property (T) groups

- $SL_n(\mathbb{Z})$, $n \geq 3$, (Kazhdan 1967), but not $SL_2(\mathbb{Z})$.
- $EL_n(\mathcal{R}) = \langle e_{ij}(r) : i \neq j, r \in \mathcal{R} \rangle \subset GL_n(\mathcal{R})$, $n \geq 3$,
where \mathcal{R} finitely generated ring and $e_{ij}(r) := I_n + rE_{ij}$
(Shalom & Vaserstein, Ershov–Jaikin–Zapirain 2006–08).
- $Aut(\mathbf{F}_n)$, $n \geq 4$. (Kaluba–Nowak–O., K–Kielak–N., Nitsche 17–20).
 $Aut(\mathbf{F}_n)$ is the noncommutative analogue of $GL_n(\mathbb{Z})$.
 $\mathbf{F}_n \twoheadrightarrow \mathbb{Z}^n$ abelianization $\rightsquigarrow Aut(\mathbf{F}_n) \twoheadrightarrow Aut(\mathbb{Z}^n) = GL_n(\mathbb{Z})$.
 $\rightsquigarrow Aut(\mathbf{F}_2)$ does not have (T). Neither $Aut(\mathbf{F}_3)$ (McCool 1989).

Product Replacement Algorithm (Celler et al. 95, Lubotzky–Pak 01)

$Aut^+(\mathbf{F}_n) = \langle R_{i,j}^\pm, L_{i,j}^\pm \rangle \leq_{\text{index } 2} Aut(\mathbf{F}_n)$, where $\mathbf{F}_n = \langle g_1, \dots, g_n \rangle$ and
 $R_{i,j}^\pm: (g_1, \dots, g_n) \mapsto (g_1, \dots, g_{i-1}, g_i g_j^\pm, g_{i+1}, \dots, g_n)$,
 $L_{i,j}^\pm: (g_1, \dots, g_n) \mapsto (g_1, \dots, g_{i-1}, g_j^\pm g_i, g_{i+1}, \dots, g_n)$.

PRA is a practical algorithm to obtain “random” elements in a given finite group Λ of rank $< n$ via the PRA random walk

$$Aut^+(\mathbf{F}_n) \curvearrowright \{(h_1, \dots, h_n) \in \Lambda^n : \Lambda = \langle h_1, \dots, h_n \rangle\}.$$

Noncommutative real algebraic geometry of property (T)

Hilbert's 17th Pb: $f \in \mathbb{R}(x_1, \dots, x_d)$, $f \geq 0$ on \mathbb{R}^d

(E. Artin 1927) $\implies f = \sum_i g_i^2$ for some $g_1, \dots, g_k \in \mathbb{R}(x_1, \dots, x_d)$.

$\mathbb{R}[\Gamma]$ real group algebra with the involution $(\sum_t \alpha_t t)^* = \sum_t \alpha_t t^{-1}$.

$\Sigma^2 \mathbb{R}[\Gamma] := \{\sum_i f_i^* f_i\} = \{\sum_{x,y} P_{x,y} x^{-1} y : P \in \mathbb{M}_\Gamma^+\}$ **positive cone**

Here \mathbb{M}_Γ^+ finitely supported positive semidefinite matrices.

- $\mathbb{B}(\mathcal{H})^+ := \{A = A^* : \langle Av, v \rangle \geq 0 \ \forall v \in \mathcal{H}\} = \Sigma^2 \mathbb{B}(\mathcal{H})$ psd operators.
- $\forall (\pi, \mathcal{H})$ unitary rep'n, $\pi(\sum_i f_i^* f_i) = \sum_i \pi(f_i)^* \pi(f_i) \geq 0$ in $\mathbb{B}(\mathcal{H})$.
- $C^*[\Gamma]$ the universal enveloping C^* -algebra of $\mathbb{R}[\Gamma]$.

Laplacian: For $\Gamma = \langle S \rangle$ with $S = S^{-1}$ finite,

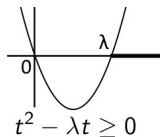
$$\Delta := \frac{1}{2} \sum_{s \in S} (1 - s)^* (1 - s) = |S| - \sum_{s \in S} s \in \Sigma^2 \mathbb{R}[\Gamma].$$

Then, $\langle \pi(\Delta)v, v \rangle = \frac{1}{2} \sum_{s \in S} \|v - \pi(s)v\|^2$ and

Γ has (T) $\iff \exists \lambda > 0 \ \forall (\pi, \mathcal{H}) \ \text{Sp}(\pi(\Delta)) \subset \{0\} \cup [\lambda, \infty)$

$\iff \exists \lambda > 0$ such that $\Delta^2 - \lambda \Delta \geq 0$ in $C^*[\Gamma]$

$$\rightsquigarrow \kappa(\Gamma, S) \geq \sqrt{2\lambda/|S|}$$



Algebraic characterization of property (T)

Let $\Gamma = \langle S \rangle$ with $S = S^{-1}$ finite.

$\mathbb{R}[\Gamma]$ real group algebra with the involution $(\sum_t \alpha_t t)^* = \sum_t \alpha_t t^{-1}$.

$$\Sigma^2 \mathbb{R}[\Gamma] := \{ \sum_i f_i^* f_i \} = \{ \sum_{x,y} P_{x,y} x^{-1} y : P \in \mathbb{M}_\Gamma^+ \}$$

Here \mathbb{M}_Γ^+ finitely supported positive semidefinite matrices.

$$\Delta := \frac{1}{2} \sum_{s \in S} (1 - s)^* (1 - s) = |S| - \sum_{s \in S} s \in \Sigma^2 \mathbb{R}[\Gamma].$$

$C^*[\Gamma]$ the universal enveloping C^* -algebra of $\mathbb{R}[\Gamma]$.

Then,

$$\begin{aligned} \Gamma \text{ has (T)} &\iff \exists \lambda > 0 \text{ such that } \Delta^2 - \lambda \Delta \geq 0 \text{ in } C^*[\Gamma] \\ &\rightsquigarrow \kappa(\Gamma, S) \geq \sqrt{2\lambda/|S|} \end{aligned}$$

Theorem (O 2013)

$$\Gamma \text{ has (T)} \iff \exists \lambda > 0 \text{ such that } \Delta^2 - \lambda \Delta \succeq 0 \text{ in } \mathbb{R}[\Gamma]$$

Stability (Netzer–Thom): It suffices if $\exists \lambda > 0 \exists \Theta \in \Sigma^2 \mathbb{R}[\Gamma]$ such that

$$\|\Delta^2 - \lambda \Delta - \Theta\|_1 \ll \lambda.$$

Semidefinite Programming (SDP)

$$\begin{aligned}\Gamma \text{ has (T)} &\iff \exists \lambda > 0 \text{ such that } \Delta^2 - \lambda \Delta \in \Sigma^2 \mathbb{R}[\Gamma] \\ &\iff \exists E \in \Gamma \exists \lambda > 0 \text{ s.t. } \Delta^2 - \lambda \Delta \in \{\sum_{x,y} P_{x,y} x^{-1} y : P \in \mathbb{M}_E^+\}\end{aligned}$$

By fixing a finite subset $E \in \Gamma$, we arrive at the SDP:

$$\begin{array}{ll}\text{minimize} & -\lambda \\ \text{subject to} & \Delta^2 - \lambda \Delta = \sum_{x,y \in E} P_{x,y} x^{-1} y, \quad P \in \mathbb{M}_E^+\end{array}$$

- Due to computer capacity limitation, we almost always take

$$E := \text{Ball}(2) = \{e\} \cup S \cup S^2 = \text{supp } \Delta \cup \text{supp } \Delta^2.$$

\rightsquigarrow Size of SDP: dimension $|E|^2$ and constraints $|E^{-1}E| = |\text{Ball}(4)|$.

Certification Procedure:

Suppose (λ_0, P_0) is a hypothetical solution obtained by a computer.

Find $P_0 \approx Q^T Q$ (with $Q\mathbf{1} = 0$) and calculate **with guaranteed accuracy**

$$r := \|\Delta^2 - \lambda_0 \Delta - \sum_{x,y} (Q^T Q)_{x,y} (1-x)^*(1-y)\|_1 \ll \lambda_0.$$

\rightsquigarrow Γ has (T) with $\lambda = \lambda_0 - 2r$ (in the case of $E = \text{Ball}(2)$).

- Solving SDP is computationally hard, but certifying (T) is relatively easy.

Results

Γ has (T) $\iff \exists E \in \Gamma \exists \lambda > 0$ s.t. $\Delta^2 - \lambda \Delta \in \{\sum_{x,y} P_{x,y} x^{-1} y : P \in \mathbb{M}_E^+\}$

Results of SDP for $E = \text{Ball}(2)$.

- $\text{SL}_n(\mathbb{Z})$ with $S = \{e_{ij}^\pm : i \neq j\}$: $\lambda_3 > 0.27$, $\lambda_4 > 1.3$, $\lambda_5 > 2.6$.
(Netzer–Thom 2014, Fujiwara–Kabaya 2017, Kaluba–Nowak 2017)
- No response for $\text{SL}_6(\mathbb{Z})$.

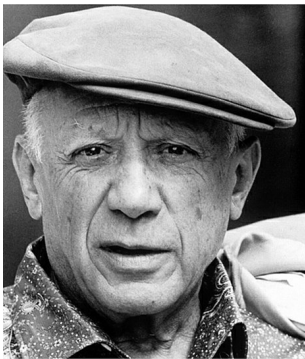
For $\text{Aut}^+(\mathbf{F}_4)$, the size of SDP $\approx 10\,000\,000$, beyond our computer's capacity. We exploited invariance under $\mathfrak{S}(n) \ltimes (\mathbb{Z}/2)^{\oplus n} \curvearrowright \text{Aut}^+(\mathbf{F}_n)$.

- $\text{Aut}^+(\mathbf{F}_4)$: ☹️☹️☹️ No response.
- $\text{Aut}^+(\mathbf{F}_5)$: !☺️^☺️^☺️! **YES!!!** with $\lambda > 1.2$.

Theorem

$\text{Aut}^+(\mathbf{F}_n)$ has property (T) for

- $n = 5$ (Kaluba–Nowak–O. 2017)
- $n \geq 6$ (Kaluba–Kielak–Nowak 2018)
- $n = 4$ (Nitsche 2020, by a new SDP)



Revista Ve a y Lea, January 1962

“But they (= computers) are useless.
They can only give you answers.”
Pablo Picasso, 1968.

Property (T) for an infinite series (KKN 2018)

$$\Gamma_n := \text{Aut}^+(\mathbf{F}_n), \quad S_n := \{R_{i,j}^\pm, L_{i,j}^\pm : i \neq j\}, \quad E_n := \{\{i,j\} : i \neq j\}$$

Want to show $\Delta_n = \sum_{s \in S_n} 1 - s$ satisfies $\Delta_n^2 - \lambda_n \Delta_n \succeq 0$.

$$\Delta_n = \sum_{e \in E_n} \Delta_e,$$

$$\begin{aligned} \Delta_n^2 &= \sum_e \Delta_e^2 + \sum_{e \sim f} \Delta_e \Delta_f + \sum_{e \perp f} \Delta_e \Delta_f \\ &=: \mathbf{Sq}_n + \mathbf{Adj}_n + \mathbf{Op}_n. \end{aligned}$$

- \mathbf{Sq}_n and \mathbf{Op}_n are positive, but \mathbf{Adj}_n may not.

For $n > m$,

$$\sum_{\sigma \in \mathfrak{S}(n)} \sigma(\Delta_m) = m(m-1) \cdot (n-2)! \cdot \Delta_n$$

$$\sum_{\sigma \in \mathfrak{S}(n)} \sigma(\mathbf{Adj}_m) = m(m-1)(m-2) \cdot (n-3)! \cdot \mathbf{Adj}_n$$

$$\sum_{\sigma \in \mathfrak{S}(n)} \sigma(\mathbf{Op}_m) = m(m-1)(m-2)(m-3) \cdot (n-4)! \cdot \mathbf{Op}_n$$

Trial and error on the computer has confirmed

$$\mathbf{Adj}_5 + \alpha \mathbf{Op}_5 - \varepsilon \Delta_5 \succeq 0$$

with $\alpha = 2$ and $\varepsilon = 0.13$. It follows that

$$0 \preceq 60(n-3)! \left(\mathbf{Adj}_n + \frac{2\alpha}{n-3} \mathbf{Op}_n - \frac{n-2}{3} \varepsilon \Delta_n \right) \preceq 60(n-3)! \left(\Delta_n^2 - \frac{n-2}{3} \varepsilon \Delta_n \right),$$

provided $2\alpha/(n-3) \leq 1$. $\rightsquigarrow \kappa(\text{Aut}^+(\mathbf{F}_n), S_n) \geq \sqrt{2\lambda_n/|S_n|} \geq \sqrt{\varepsilon/6n}$

Generalizing property (T) for $\text{EL}_n(\mathcal{R})$ for a rng \mathcal{R}

The computer taught us the inequality $\mathbf{Adj}_5 + \alpha \mathbf{Op}_5 - \varepsilon \Delta_5 \succeq 0$ is true, useful, and even **easy to prove (!) when $\alpha > 0$ is large**.

Theorem (O. 2022)

For any f.g. **comm.** rng \mathcal{R} generated by $R_0 \subseteq \mathcal{R}$ and for n **large enough**,

$$\Delta := \sum_{r \in R_0} \sum_{i \neq j} (1 - e_{ij}(r))^* (1 - e_{ij}(r)) \quad \text{and}$$

$$\Delta^{(2)} := \sum_{r, s \in R_0} \sum_{i \neq j} (1 - e_{ij}(rs))^* (1 - e_{ij}(rs))$$

in $\mathbb{R}[\text{EL}_n(\mathcal{R})]$ satisfy

$$\Delta^2 \geq \varepsilon \Delta^{(2)}$$

in $C^*[\text{EL}_n(\mathcal{R})]$ for some $\varepsilon > 0$ (but not $\Delta^2 \succeq \varepsilon \Delta^{(2)}$ in $\mathbb{R}[\text{EL}_n(\mathcal{R})]$).

⚠ “rng” = “ring” – “i”. $\text{EL}_n(\mathcal{R}) \twoheadrightarrow \text{EL}_n(\mathcal{R}/\mathcal{R}^2) \cong (\mathcal{R}/\mathcal{R}^2)^{\oplus n(n-1)}$ abelian.

The proof of $\mathbf{Adj}_5 + \alpha \mathbf{Op}_5 - \varepsilon \Delta_5^{(2)} \succeq 0$ is silicon-free and relies on Boca and Zaharescu’s work (2005) on the almost Mathieu operators in \mathcal{A}_θ .

Corollary

$\exists n \exists \varepsilon > 0$ s.t. $\text{Cayley}(\text{SL}_n(\mathbb{Z}/q\mathbb{Z}), \{e_{ij}(p) : i \neq j\})$, $p \perp q$, are ε -expanders.

⚠ The Kazhdan constants for $(\text{EL}_n(p\mathbb{Z}))_{p \in \mathbb{N}}$ are not uniform.